

Logical Loops

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A general iteration construct for Prolog and Prolog-based (CLP) languages

Definition via program transformation

Comparison with alternatives

Recursion Higher-order constructs Bounded Quantifiers

Expressive Power

Experience



Introduced in ECLiPSe to support

Prolog beginners

Provide a familiar feature: loops

Application writers

Improved productivity, code structure, maintainability

Constraint Problem Modellers

Provide equivalent of bounded quantification



Iteration (over list elements)

Replace this

With this

).

```
write_list(List) :-
```

```
write("List: "),
```

```
write_list1(List).
```

```
write_list(List) :-
```

```
write("List: "),
```

```
( foreach(X,List) do
    write(X)
```

```
write_list1([]).
write_list1([X|T]) :-
    write(X),
    write_list1(T).
```

```
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```

Iteration (over range of numbers)

write_nat1(I, N).

Replace thisWith this

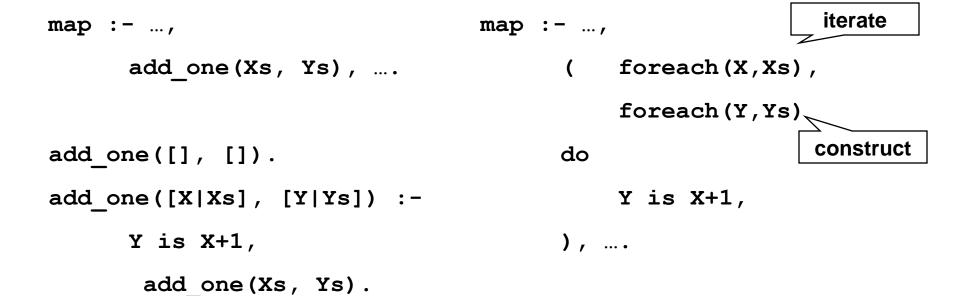
Replace this





Replace this







Programming idioms covered

Fully:

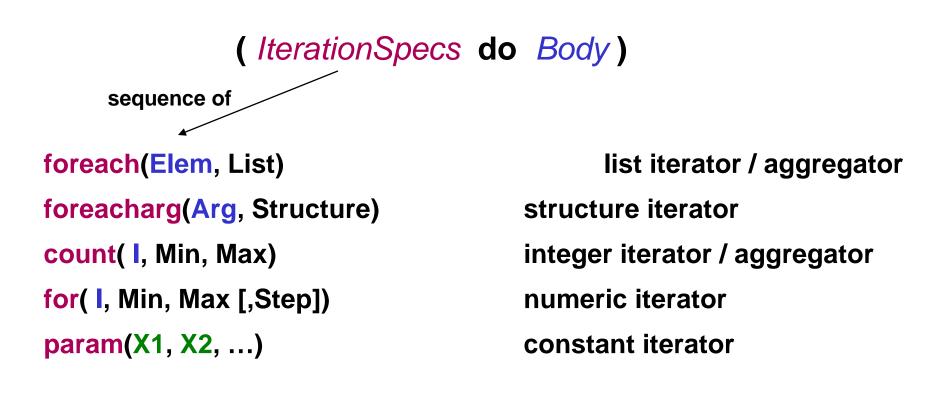
- Iteration
- Aggregation
- Mapping

Partly:



While-loop

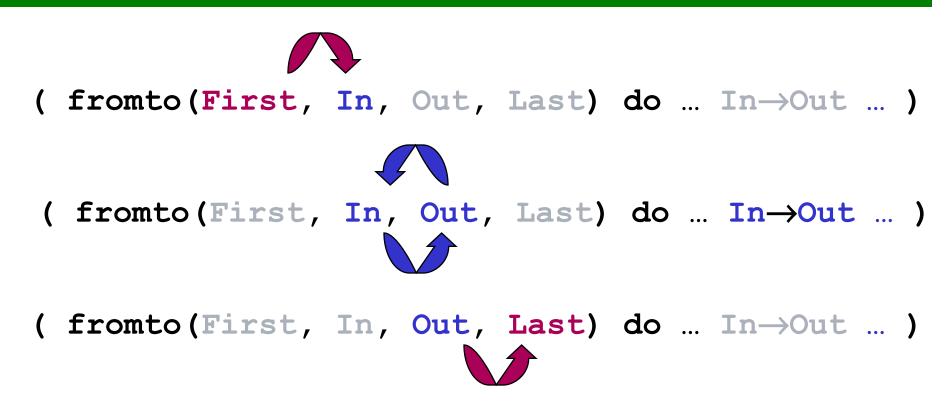




fromto(First, In, Out, Last)

generic iterator / aggregator

The generic fromto-iterator



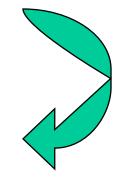
Fromto can express all other iterators, e.g.

fromto(List, [X|Xs], Xs, []) ⇔ foreach(X,List)

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A goal

..., (IterationSpecifiers do Body), ... Is replaced by ..., PreCallGoals, µ(CallArgs), ... With an auxiliary predicate µ defined as μ(**BaseArgs**) :- !. μ(HeadArgs) :-PreBodyGoals, Body, μ(*RecArgs*).





% Initial call

% PreCallGoals

- % Base clause
- % Recursive clause head
- % PreBodyGoals
- % Body
- % Recursive call

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Loops vs. Recursion

Conciseness

Auxiliary predicate and its arity are hidden.

Modifiability

Add/remove single iteration specifier rather than argument(s) in 4 places.

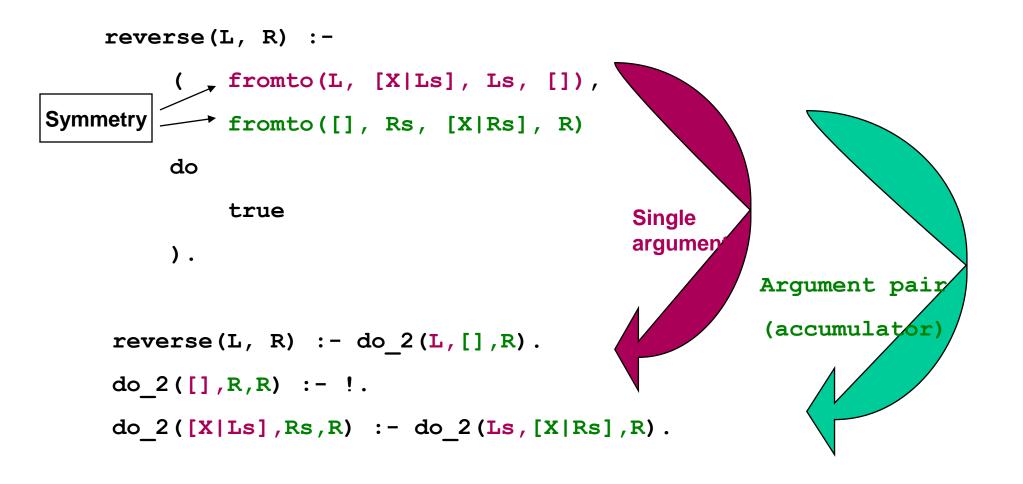
Structure

Nested loops can be clearer than a flat collection of predicates. Iteration specifiers group related information together.

Abstraction

Intention of iteration is explicit.

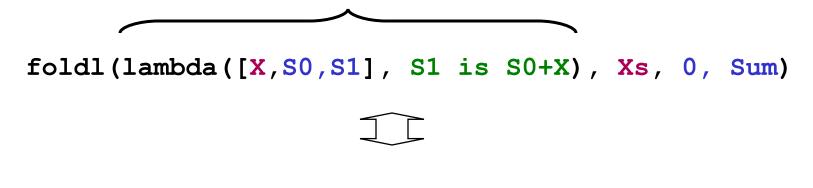
Fromto can help beginners understand the concept of accumulators.



Higher-Order constructs for similar tasks:

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More realistic comparison with a lambda-term syntax:



foreach(X,Xs), fromto(0,S0,S1,Sum) do S1 is S0+X

Loop formulation has same size and better grouping

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Loops vs Higher-Order Constructs (3)

Higher-Order map/fold/filter

- Data structure (list) specific
- Arbitrary traversal orders
- foldl / foldr symmetry
- Families needed: map/(2+N), fold/(2+2N), filter/(2+N)
- Combinations needed: map_foldl/5, …
- Higher-order-typed arguments

Do-loop

- Any data structure (fromto)
- Only iteration / tail recursion
- Only efficient foldl-equivalent
- Arbitrary number of iteration specifiers allowed
- Arbitrary combinations of iteration specifiers allowed
- Same typing as recursive form

Claim: Higher-order constructs are powerful, but not the best for iteration

Quantification over finite sets

Mathematical modelling languages (AMPL and others) Extension to LP (Voronkov 90, Barklund et al 93/95, Apt 96) $pos_list(List) := \forall X \in List X > 0.$ $pos_array(A,N) := \forall I \in 1..N A[I] > 0.$

But: mapping of lists is impossible

add_one_to_all(Xs,Ys) :- $\forall x \in Xs \ \forall Y \in Ys \ Y \ is \ X+1$. Does not work: No concept of "corresponding elements" Needs arrays to be useful !

But: aggregation operators

E.g. minimum, maximum, sum, ... must be introduced separately
 sum_list(Xs,S) :- S=ΣX : X∈Xs.

Expressive Power of Loops (1)

(

)

• We can write a while-loop, albeit awkwardly:

```
fromto(cont, _, Continue, stop)
do
  ( termination_condition(...) ->
      Continue = stop
  ; Continue = cont
  ),
...
```



A recursion-free meta-interpreter for pure Prolog:

```
solve(Q) :-
   ( fromto([Q], [G|C0], C1, []) do
        solve_step(G, C0, C1)
   ).
   solve_step(true, C, C).
   solve_step((A,B), C, [A,B|C]).
   solve step(A, C, [B|C]) :- clause(A, B).
```

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A language feature should

- Fit with existing language concepts / user ideas "Principle of least astonishment"
- Provide a clear advantage Code size, elegance, maintainability, robustness, ...

Not have an overhead cost

Otherwise programmers will use lower-level methods

Loops have been in ECLiPSe officially since 1998.

Analysis of commerical code developed with ECLiPSe:

- Application A (1997-2001)
 254 predicate in 24 modules
 34 loops
- Application B (2000-2002)
 528 predicates in 35 modules
 210 loops



Language extension with a minimum of new concepts no types, modes, arrays were needed

 Efficiently implementable virtually no runtime overheads

Accepted by programmers

this has always been a problem with higher-order constructs

Translation code available at

http://www.icparc.ic.ac.uk/eclipse/software/loops/

Nondeterministic iteration

Once the termination conditions are met, the loop terminates. Allowing more iterations on backtracking seems too error-prone.

Better support for while/until-loops

Makes the semantics much more complex.

Clearly separate iterators and aggregators Would lose some multi-directionality, but a good idea with modes.

More/less iterator shorthands

A matter of taste.

Array notation makes loops even more useful:

```
...,
( for(I,1,N), param(Board,N) do
        ( for(J,I+1,N), param(Board,I) do
        Board[I] #\= Board[J],
        Board[I] #\= Board[J]+J-I,
        Board[I] #\= Board[J]+I-J
        )
).
```

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Bounded Quantification

• "A priori" Bounded Quantification

(Barklund and Bevemyr (Reform Prolog) 1993) (Barklund and Hill (Gödel) 1995) (Apt 1996) Mathematical modelling languages

Quantification over finite sets:

pos_list(List) :- $\forall X \in List$ X > 0. pos_array(A,N) :- $\forall I \in 1..N$ A[I] > 0.

Bounded Quantification (Voronkov 1990)

More powerful due to list-suffix quantifier Termination can depend on quantified formula Turing-complete



Mapping of lists is impossible with a priori BQ

add_one_to_all(Xs,Ys) :- $\forall x \in Xs \ \forall Y \in Ys \ Y \ is \ X+1$. Does not work: No concept of "corresponding elements"

A priori BQ needs arrays

add_one_to_all(A,B) :- $\forall I \in 1..N$ B[I] is A[I]+1. Indeed all authors introduce arrays into their proposed languages. But still cannot map a list to an array ...

Voronkov's full BQ can express list mapping However, very unintuitive

Loop solution:

We have a concept of implicitly ordered iteration steps

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Aggregation not covered by basic a priori BQ

 Specific aggregation operators must be introduced e.g. minimum, maximum, sum, ... sum_list(Xs,S) :- S=ΣX : X∈Xs.

Loop solution:

The fromto-specifier can function as iterator and aggregator. Any aggregation function can be expressed.

